

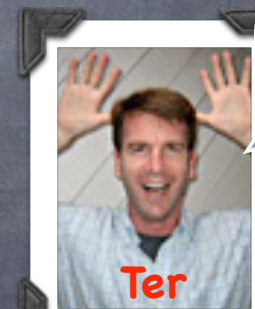
ANTLR 3

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2008

ANTLR Overview



- ANother Tool for Language Recognition
 - written by Terence Parr in Java
- Easier to use than most/all similar tools
- Supported by ANTLRWorks
 - graphical grammar editor and debugger
 - written by Jean Bovet using Swing
- Used to implement
 - “real” programming languages
 - domain-specific languages (DSLs)
- <http://www.antlr.org>
 - download ANTLR and ANTLRWorks here
 - both are free and open source
 - docs, articles, wiki, mailing list, examples



Ter

I'm a professor at the University of San Francisco.



Jean

I worked with Ter as a masters student there.

ANTLR Documentation

<http://antlr.org>

[Home](#) | [Download](#) | [ANTLRWorks](#) | [Wiki](#) | [About ANTLR](#) | [Feedback](#) | [Support](#) | [Bugs](#) | [v2](#)

ANTLR v3

Latest version is 3.0.0
Download now! [DOWNLOAD](#)

What is ANTLR?

ANTLR, ANOther Tool for Language Recognition, is a language tool that provides a framework for constructing recognizers, interpreters, compilers, and translators from grammatical descriptions containing actions in a variety of [target languages](#). ANTLR provides excellent support for tree construction, tree walking, translation, error recovery, and error reporting. There are currently about 5,000 ANTLR source downloads a month.



ANTLR has a sophisticated grammar development environment called [ANTLRWorks](#), written by Jean Devrieux.



[Terence Parr](#) is the maniac behind ANTLR and has been working on language tools since 1989. He is a professor of computer science at the [University of San Francisco](#).

[More...](#)

Testimonials

Turning PHP into a functional PL

Jeffrey M. Barber

Antlr v3 is awesome. I used Antlr v2 for several projects, but my latest...

You Used Ruby to Write WHAT?!

Zed Shaw

"...using ANTLR, without much fuss I can prototype an entire new language..."

Regarding The Definitive ANTLR Reference book

Gevik Babakhani

Before I got this book, I had to hack my way through various examples and...

Still using ANTLR after all these years

Ron Ten-Hove

I've been using ANTLR since the first SIGPLAN Notices printing of the PCCTS...

[More...](#)

Showcase

[New version of ANTLR Tester](#)

Jeremy D. Frens *Thu Mar 27, 2008 10:33*

The ANTLR Testing library is a JUnit extension to test ANTLR grammars....

[Adobe Flex Builder 3](#)

Steve Breinberg *Wed Mar 26, 2008 16:45*

Adobe(r) Flex(r) Builder(tm) 3 software is a highly productive Eclipse(tm)...

[SelectView](#)

Alien~ *Sun Feb 17, 2008 23:05*

A tool for Relational Data Analysis. split show relational data.

[More...](#)

Looking for previous version [ANTLR v2?](#)

If you like ANTLR, check out the [StringTemplate template engine](#).

SEARCH

News

News feed has moved to [wiki](#).

[ANTLR news...](#)

[Terence's blog...](#)

File Sharing

[Sun-tuned ANTLR v2](#)

Sun Microsystems / NetBeans *Tue Jun 3, 2008 14:22*

Documentation

[Getting started with ANTLR v3](#)

[ANTLR Documentation](#)

[The Definitive ANTLR Reference: Building domain-specific languages](#)

Terence's ANTLR v3 book is now available (May, 2007). You can buy the PDF of it also.

[ANTLR API Documentation](#)

Articles

[ANTLR 3.0](#)

Mark Volkmann *Mon Jun 2, 2008 12:18*
A large article talking about how to use ANTLR 3.0.

[The Reuse of Grammars with Embedded Semantic Actions](#)

Terence Parr *Thu Apr 3, 2008 10:33*
My keynote paper for IEEE International Conference on Program Comprehension...

[Create Domain-Specific Languages with ANTLR](#)

Rod Coffin and Paul Holser *Wed Nov 14, 2007 14:45*



ANTLR Overview ...

- Uses EBNF grammars

- Extended Backus-Naur Form
- can directly express optional and repeated elements
- supports subrules (parenthesized groups of elements)

BNF grammars require more verbose syntax to express these.

- Supports many target languages for generated code

- Java, Ruby, Python, Objective-C, C, C++ and C#

- Provides infinite lookahead

- most parser generators don't
- used to choose between rule alternatives

- Plug-ins available for IDEA and Eclipse

ANTLR Overview ...

• Three main use cases

We'll explain actions and rewrite rules later.

• 1) Implementing "validators"

no actions or rewrite rules

- generate code that validates that input obeys grammar rules

• 2) Implementing "processors"

actions but no rewrite rules

- generate code that validates and processes input
- could include performing calculations, updating databases, reading configuration files into runtime data structures, ...
- our Math example coming up does this

• 3) Implementing "translators"

actions containing printlns and/or rewrite rules

- generate code that validates and translates input into another format such as a programming language or bytecode

Projects Using ANTLR

Programming languages

Boo

- <http://boo.codehaus.org>

Groovy

- <http://groovy.codehaus.org>

Mantra

- <http://www.linguamantra.org>

Nemerle

- <http://nemerle.org>

XRuby

- <http://xruby.com>

Other tools

Hibernate

- for its HQL to SQL query translator

IntelliJ IDEA

Jazillian

- translates COBOL, C and C++ to Java

JBoss Rules (was Drools)

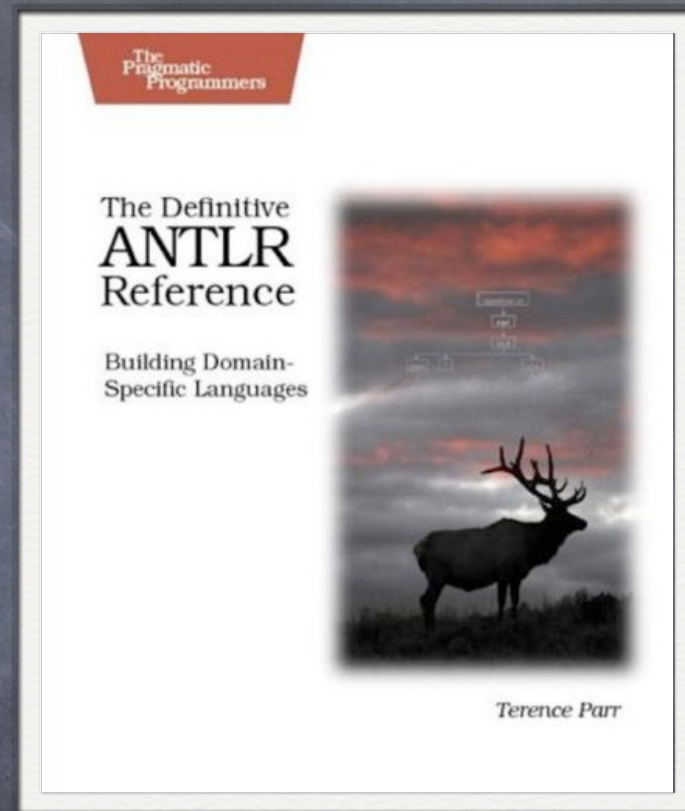
Keynote (Apple)

WebLogic (Oracle)

too many more list!

See showcase and testimonials at
<http://antlr.org/showcase/list> and
<http://www.antlr.org/testimonial/>.

Books



- “ANTLR Recipes”? in the works
 - another Pragmatic Programmers book from Terence Parr

Other DSL Approaches

- Languages like Ruby and Groovy are good at implementing DSLs, but ...
- The DSLs have to live within the syntax rules of the language
- For example
 - dots between object references and method names
 - parameters separated by commas
 - blocks of code surrounded by { ... } or do ... end
- What if you don't want these in your language?

Conventions

- ANTLR grammar syntax makes frequent use of the characters [] and { }
- In these slides
 - when describing a placeholder, I'll use *italics*
 - when describing something that's optional, I'll use *item?*

Some Definitions

• Lexer

- converts a stream of characters to a stream of tokens

• Parser

Token objects know their start/stop character stream index, line number, index within the line, and more.

- processes a stream of tokens, possibly creating an AST

• Abstract Syntax Tree (AST)

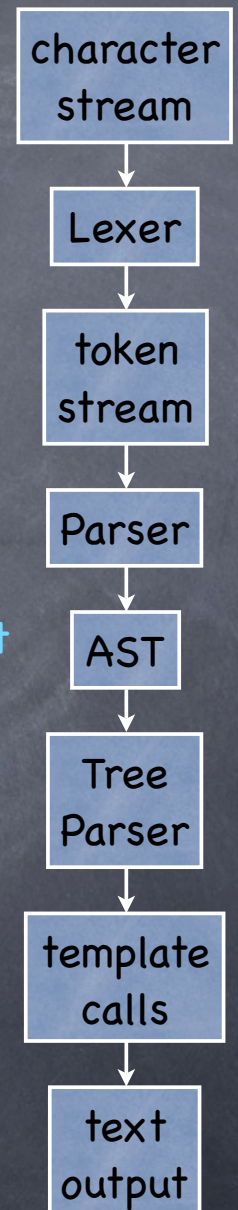
- an intermediate tree representation of the parsed input that
 - is simpler to process than the stream of tokens
 - can be efficiently processed multiple times

• Tree Parser

- processes an AST

• StringTemplate

- a library that supports using templates with placeholders for outputting text (for example, Java source code)



General Steps

- Write grammar
 - can be in one or more files
- Optionally write StringTemplate templates
- Debug grammar with ANTLRWorks
- Generate classes from grammar
 - these validate that text input conforms to the grammar and execute target language "actions" specified in the grammar
- Write application that uses generated classes
- Feed the application text that conforms to the grammar

Let's Create A Language!

Features

- run on a file or interactively
- get help - ? or help
- one data type, double
- assign values to variables - $a = 3.14$
- define polynomial functions - $f(x) = 3x^2 - 4x + 2$
- print strings, numbers, variables and function evaluations -
`print "The value of f for " a " is " f(a)`
- print the definition of a function and its derivative -
`print "The derivative of " f() " is " f'()`
- list variables and functions -
`list variables and list functions`
- add/subtract functions - $h = f - g$ ←
 - the function variables don't have to match
- exit - exit or quit

Input:

```
f(x) = 3x^2 - 4
g(y) = y^2 - 2y + 1
h = f - g
print h()
```

Output:

```
h(x) = 2x^2 + 2x - 5
```

Example Input/Output

```
a = 3.14
f(x) = 3x^2 - 4x + 2
print "The value of f for " a " is " f(a)

print "The derivative of " f() " is " f'()
```

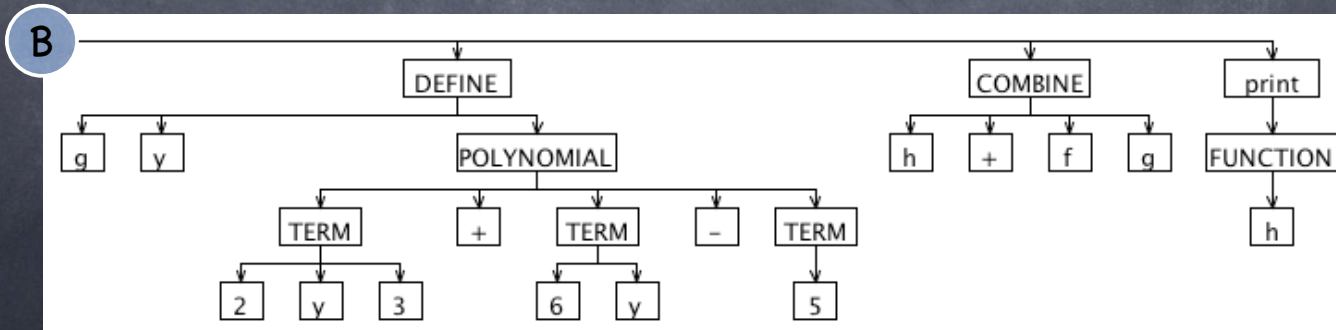
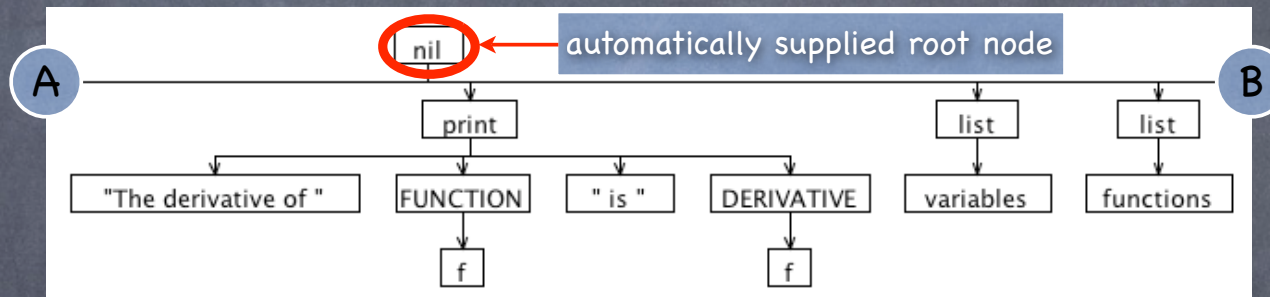
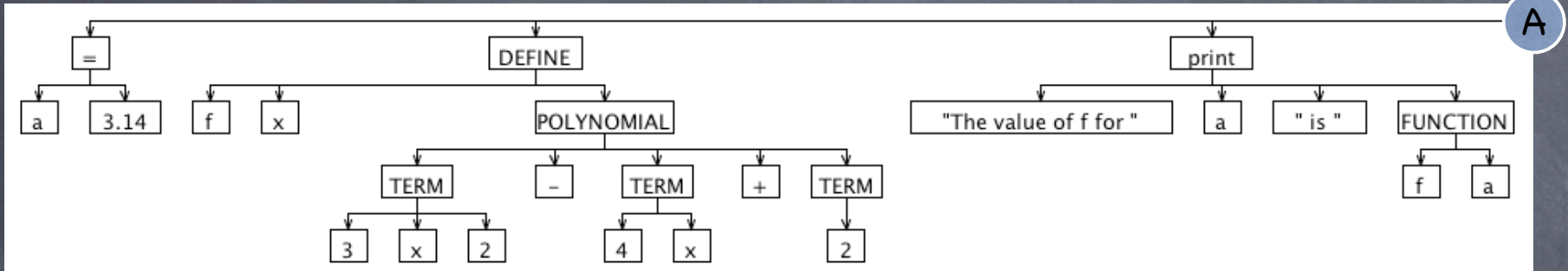
```
list variables
list functions
```

```
g(y) = 2y^3 + 6y - 5
h = f + g
print h()
```

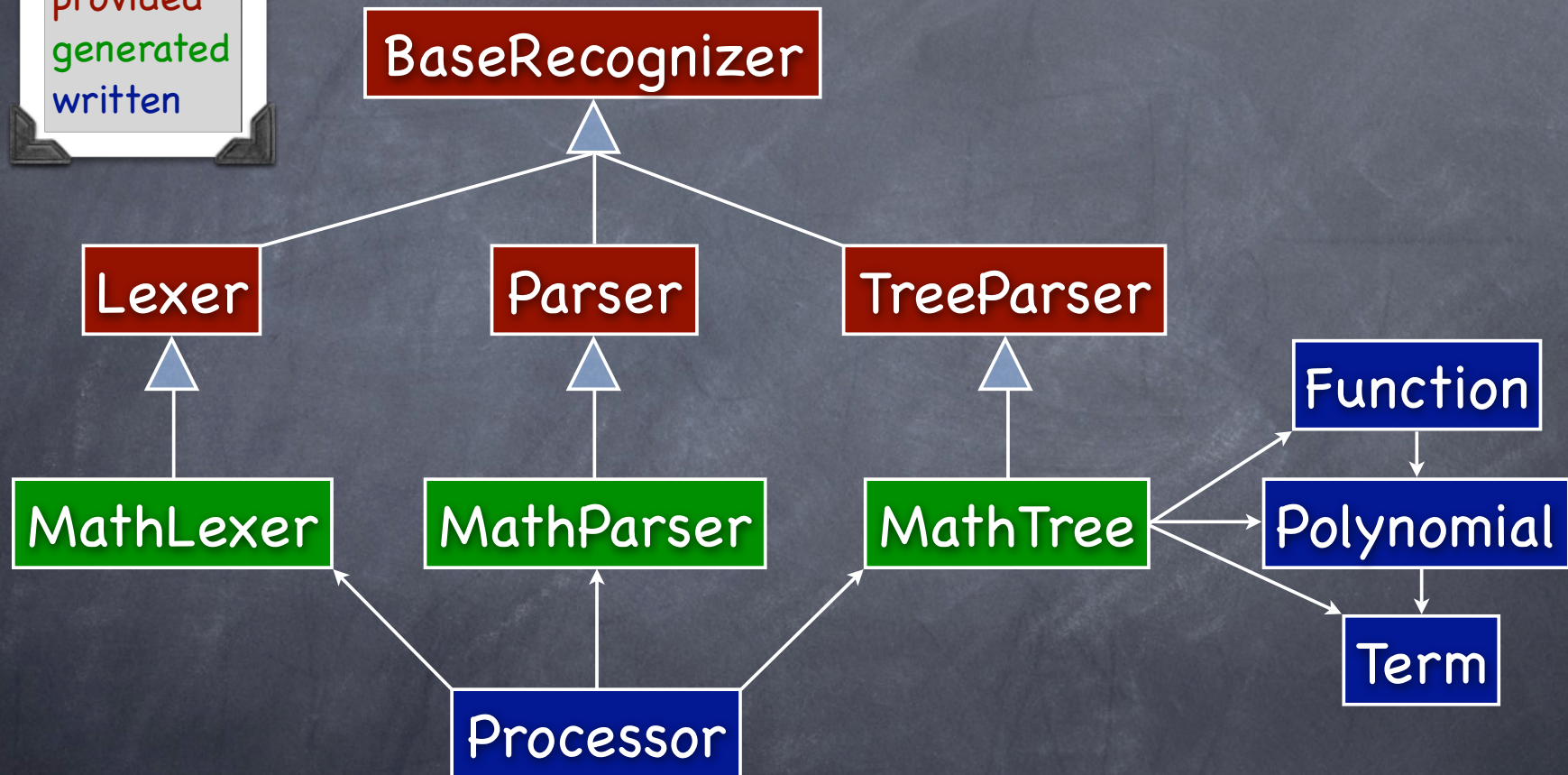
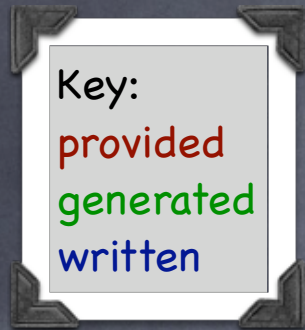
```
The value of f for 3.14 is 19.0188
The derivative of f(x) = 3x^2 - 4x + 2
is f'(x) = 6x - 4
# of variables defined: 1
a = 3.14
# of functions defined: 1
f(x) = 3x^2 - 4x + 2
h(x) = 2x^3 + 3x^2 + 2x - 3
```

Example AST

drawn by
ANTLRWorks



Important Classes



Grammar Actions

- Add to the generated code
- `@grammar-type::header { ... }`

- inserts contained code before the class definition
 - commonly used to specify a package name and import classes in other packages

grammar-type
must be lexer,
parser (the default)
or treeparser

- `@grammar-type::members { ... }`

- inserts field declarations and methods inside the class definition
- commonly used to
 - define constants and attributes accessible to all rule methods in the generated class
 - define methods used by multiple rule actions
 - override methods in the superclasses of the generated classes
 - useful for customizing error reporting and handling

Lexer Rules

- Need one for every kind of token to be processed in parser grammar
- Name must start uppercase
 - typically all uppercase
- Assign a token name to
 - a single literal string found in input
 - a selection of literal strings found in input
 - one or more characters and ranges of characters
 - can use cardinality indicators ?, * and +
- Can refer to other lexer rules
- “fragment” lexer rules
 - do not result in tokens
 - are only referenced by other lexer rules



Regular expressions aren't supported.

The next lexer rule used is the one that matches the most characters. If there is a tie, the one listed first is used, so order matters!

See LETTER and DIGIT rules in the upcoming example.

Whitespace & Comments

- Handled in lexer rules
- Two common options
 - throw away - `skip()`;
 - write to a different "channel" - `$channel = HIDDEN;`

The ANTLRWorks debugger input panel doesn't display skipped characters, but does display hidden ones.

constant defined in BaseRecognizer

Examples

```
WHITESPACE: (' ' | '\t')+ { $channel = HIDDEN; };
NEWLINE: ('\r'? '\n')+;
SINGLE_COMMENT: '//' ~('\r' | '\n')* NEWLINE { skip(); };
MULTI_COMMENT
options { greedy = false; }
: '/' '*' .* '/' NEWLINE? { skip(); };
```

Don't skip or hide NEWLINES if they are used as statement terminators.

The greedy option defaults to true, except for the patterns `.*` and `.+`, so it doesn't need to be specified here. When true, the lexer matches as much input as possible. When false, it stops when input matches the next element.

Our Lexer Grammar

```
lexer grammar MathLexer;
```

```
@header { package com.ociweb.math; }
```

We want the generated lexer class to be in this package.

```
APOSTROPHE: '\''; // for derivative
```

```
ASSIGN: '=';
```

```
CARET: '^'; // for exponentiation
```

```
FUNCTIONS: 'functions'; // for list command
```

```
HELP: '?' | 'help';
```

```
LEFT_PAREN: '(';
```

```
LIST: 'list';
```

```
PRINT: 'print';
```

```
RIGHT_PAREN: ')';
```

```
SIGN: '+' | '-';
```

```
VARIABLES: 'variables'; // for list command
```

```
NUMBER: INTEGER | FLOAT;
```

```
fragment FLOAT: INTEGER '.' '0'..'9'+;
```

```
fragment INTEGER: '0' | SIGN? '1'..'9' '0'..'9'*;
```

Our Lexer Grammar ...

```
NAME: LETTER (LETTER | DIGIT | '_' )*;
STRING_LITERAL: '"' NONCONTROL_CHAR* '"';

fragment NONCONTROL_CHAR: LETTER | DIGIT | SYMBOL | SPACE;
fragment LETTER: LOWER | UPPER;
fragment LOWER: 'a'..'z';
fragment UPPER: 'A'..'Z';
fragment DIGIT: '0'..'9';
fragment SPACE: ' ' | '\t';

// Note that SYMBOL omits the double-quote character,
// digits, uppercase letters and lowercase letters.
fragment SYMBOL: '!' | '#'..'/' | ':'..'@' | '['..'\' | '{'..'~';

// Windows uses \r\n. UNIX and Mac OS X use \n.
// To use newlines as a terminator,
// they can't be written to the hidden channel!
NEWLINE: ('\r'? '\n')+;
WHITESPACE: SPACE+ { $channel = HIDDEN; };
```

Token Specification

- The lexer creates tokens for all input character sequences that match lexer rules
- It can be useful to create other tokens that

- don't exist in the input (imaginary)

See all the uppercase token names in the AST diagram on slide 14.

- often serve to group other tokens

- have a better name than is found in the input

- Do this with a token specification in the parser grammar

We need this for the imaginary tokens DEFINE, POLYNOMIAL, TERM, FUNCTION, DERIVATIVE and COMBINE.

```
tokens {  
    imaginary-name;  
    better-name = 'input-name';  
    . . .  
}
```

Rule Syntax

only for
lexer rules



fragment? *rule-name* *arguments?*

(returns return-values)?

throws-spec?

rule-options?

rule-attribute-scopes?

→ *rule-actions?*

: token-sequence-1

| token-sequence-2

...

;

exceptions-spec?

to customize exception
handling for this rule

add code before
and/or after code in
the generated method
for this rule

Each element in these alternative sequences
can be followed by an action which is
target language code in curly braces.
The code is executed immediately after
a preceding element is matched by input.

Creating ASTs

- Requires grammar option `output = AST;`

- Approach #1 – Rewrite rules

- appear after a rule alternative
- the recommended approach in most cases
- `-> ^(parent child-1 child-2 ... child-n)`

can't use both approaches in the same rule alternative!

- Approach #2 – AST operators

- appear in a rule alternative, immediately after tokens
- works best for sequences like mathematical expressions
- operators

- `^` – make new root node for all child nodes at the same level
- `none` – make a child node of current root node
- `!` – don't create a node ←

often used for bits of syntax that aren't needed in the AST such as parentheses, commas and semicolons

- `parent^ '(! child-1 child-2 ... child-n ')!`

Declaring Rule Arguments and Return Values

```
rule-name[type1 name1, type2 name2, ...]  
returns [type1 name1, type2 name2, ...] :  
    ...  
;
```

arguments

return values;
can have more than one

ANTLR generates a class to use as the return type of the generated method for the rule.

Instances of this class hold all the return values.

The generated method name matches the rule name.

The name of the generated return type class is the rule name with “_return” appended.

Our Parser Grammar

```
parser grammar MathParser;
```

```
options {
```

```
    output = AST;
```

```
    tokenVocab = MathLexer;
```

```
}
```

We're going to output an AST.

We're going to use the tokens defined in our MathLexer grammar.

```
tokens {
```

```
    COMBINE;
```

```
    DEFINE;
```

```
    DERIVATIVE;
```

```
    FUNCTION;
```

```
    POLYNOMIAL;
```

```
    TERM;
```

```
}
```

These are imaginary tokens that will serve as parent nodes for grouping other tokens in our AST.

```
@header { package com.ociweb.math; }
```

We want the generated parser class to be in this package.

Our Parser Grammar ...

```
// This is the "start rule".
```

```
script: statement* EOF;
```

AST operator

EOF is a predefined token that represents the end of input. The start rule should end with this.

```
statement: assign | define | interactiveStatement | combine | print;
```

```
interactiveStatement: help | list;
```

An expression starting with ">" is called a "rewrite rule".

```
assign: NAME ASSIGN value terminator -> ^(ASSIGN NAME value);
```

Examples:
a = 19
a = b
a = f(2)
a = f(b)

Parts of rule alternatives can be assigned to variables (ex. `fn` & `v`) that are used to refer to them in rule actions. Alternatively rule names (ex. `NAME`) can be used.

```
value: NUMBER | NAME | functionEval;
```

```
functionEval
```

```
: fn=NAME LEFT_PAREN (v=NUMBER | v=NAME) RIGHT_PAREN -> ^(FUNCTION $fn $v);
```

Examples:
f(2)
f(b)

```
// EOF cannot be used in lexer rules, so we made this a parser rule.
```

```
// EOF is needed here for interactive mode where each line entered ends in EOF
```

```
// and for file mode where the last line ends in EOF.
```

```
terminator: NEWLINE | EOF;
```

When parser rule alternatives contain literal strings, they are converted to references to automatically generated lexer rules. For example, we could eliminate the `ASSIGN` lexer rule and change `ASSIGN` to `'='` in this grammar. The rules in this grammar don't use literal strings.

Our Parser Grammar ...

define

```
: fn=NAME LEFT_PAREN fv=NAME RIGHT_PAREN ASSIGN  
  polynomial[$fn.text, $fv.text] terminator  
-> ^(DEFINE $fn $fv polynomial);
```

// fnt = function name text; fvt = function variable text

polynomial[String fnt, String fvt]

```
: term[$fnt, $fvt] (SIGN term[$fnt, $fvt])*  
-> ^(POLYNOMIAL term (SIGN term)*);
```

Examples:

$f(x) = 3x^2 - 4$

$g(y) = y^2 - 2y + 1$

To get the text value from a variable that refers to a Token object, use "\$var.text".

Examples:

$3x^2 - 4$

$y^2 - 2y + 1$

Our Parser Grammar ...

```
// fnt = function name text; fvt = function variable text
```

```
term[String fnt, String fvt]
```

```
// tv = term variable
```

```
: c=coefficient? (tv=NAME e=exponent?)? ←
```

```
// What follows is a validating semantic predicate.
```

```
// If it evaluates to false, a FailedPredicateException will be thrown.
```

```
{ tv == null ? true : ($tv.text).equals($fvt) }?
```

```
-> ^(TERM $c? $tv? $e?)
```

```
;
```

```
catch [FailedPredicateException fpe] {
```

```
String tvt = $tv.text;
```

```
String msg = "In function \"" + fnt +
```

```
  "\" the term variable \"" + tvt +
```

```
  "\" doesn't match function variable \"" + fvt + "\".";
```

```
throw new RuntimeException(msg);
```

```
}
```

```
coefficient: NUMBER;
```

```
exponent: CARET NUMBER -> NUMBER; ←
```

Examples:

4

4x

x^2

4x^2

Term variables must match their function variable. This catches bad function definitions such as $f(x) = 2y$.

Example:
^2

Our Parser Grammar ...

help: HELP terminator -> HELP;

Examples:
?
help

list

: LIST listOption terminator -> ^(LIST listOption);

Examples:
list functions
list variables

listOption: FUNCTIONS | VARIABLES;

Examples:
functions
variables

combine

: fn1=NAME ASSIGN fn2=NAME op=SIGN fn3=NAME terminator
-> ^(COMBINE \$fn1 \$op \$fn2 \$fn3);

Examples:
h = f + g
h = f - g

Our Parser Grammar ...

print

: PRINT printTarget* terminator -> ^(PRINT printTarget*);

Example:

print "f(" a ") = " f(a)

printTarget

: NUMBER -> NUMBER

| s1=STRING_LITERAL -> \$s1

| NAME -> NAME

// This is a function reference to print a string representation.

| NAME LEFT_PAREN RIGHT_PAREN -> ^(FUNCTION NAME)

| functionEval

| derivative

;

Examples:

19

3.14

"my text"

a

f()

f(2)

f(a)

f'()

derivative

: NAME APOSTROPHE LEFT_PAREN RIGHT_PAREN -> ^(DERIVATIVE NAME);

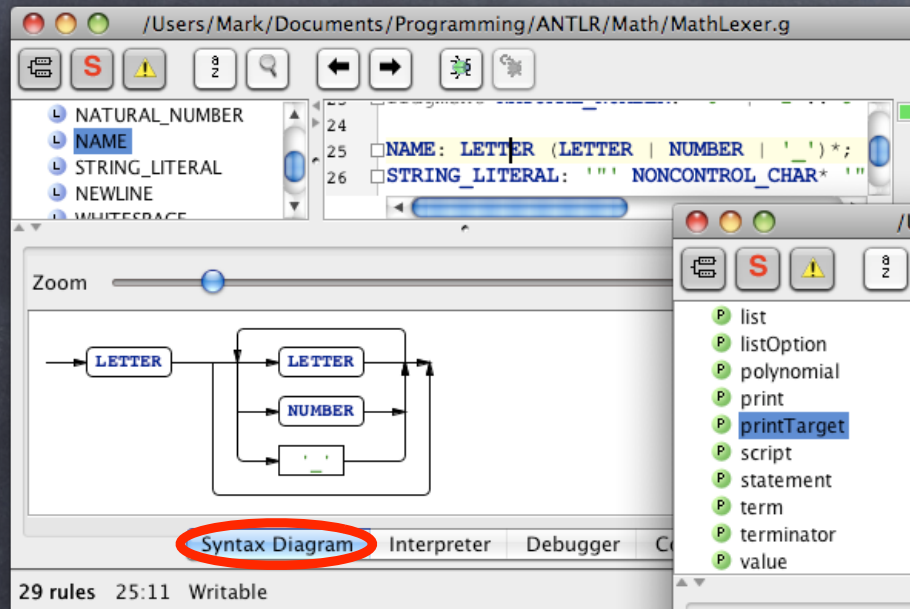
Example:

f'()

ANTLRWorks

- A graphical grammar editor and debugger
- Features
 - highlights grammar syntax errors
 - checks for grammar errors beyond the syntax variety
 - such as conflicting rule alternatives
 - displays a syntax diagram for the selected rule
 - debugger can step through creation of parse trees and ASTs

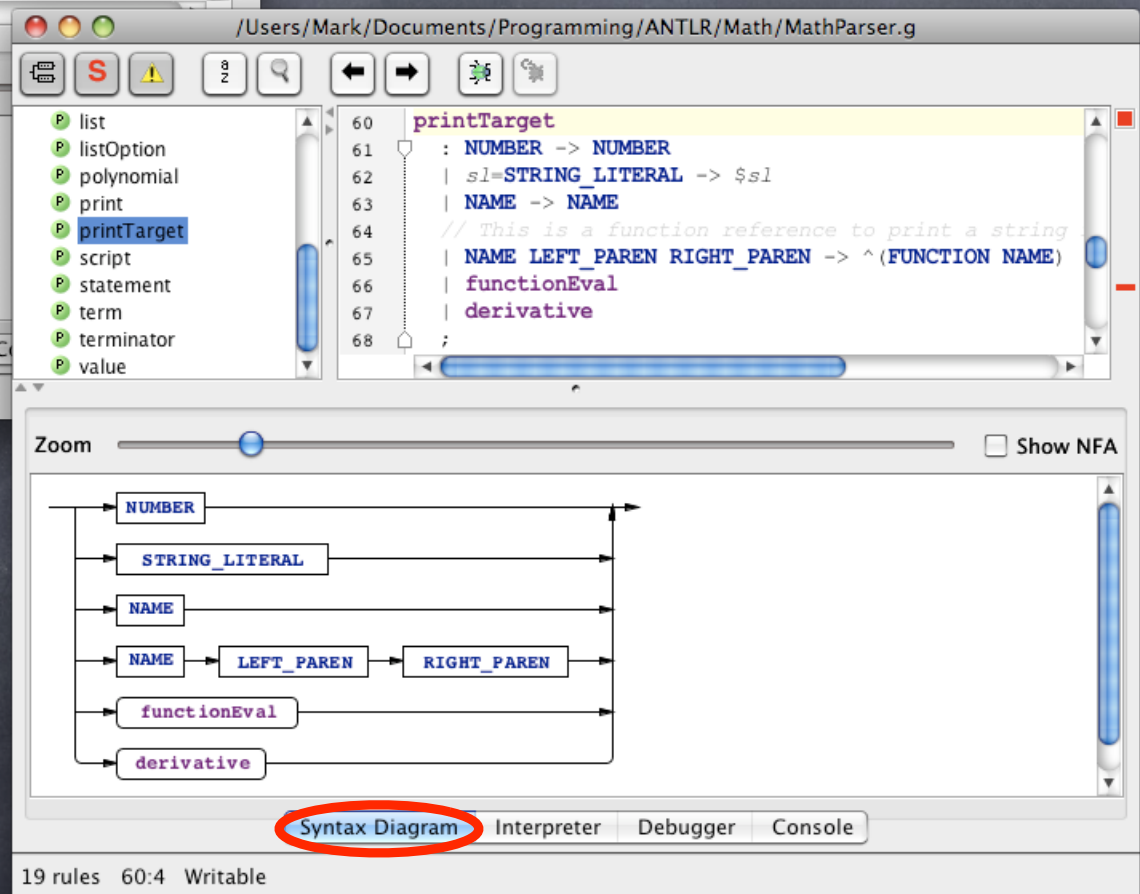
ANTLRWorks ...



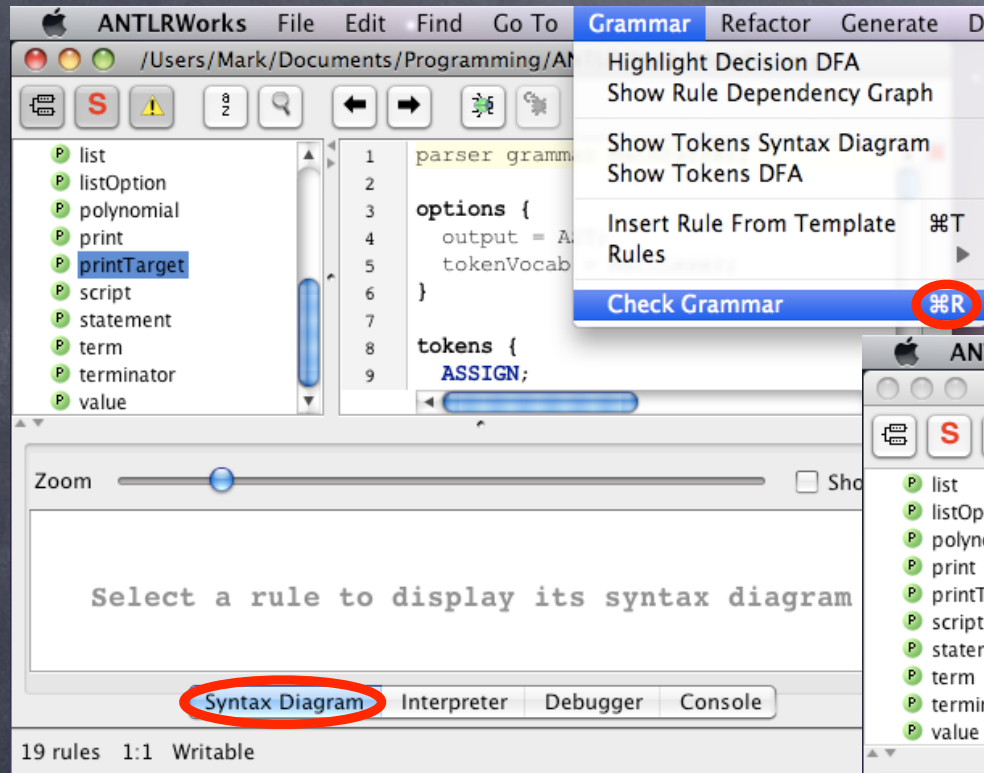
lexer rule
syntax diagram

Rectangles correspond
to fixed vocabulary symbols.
Rounded rectangles correspond
to variable symbols.

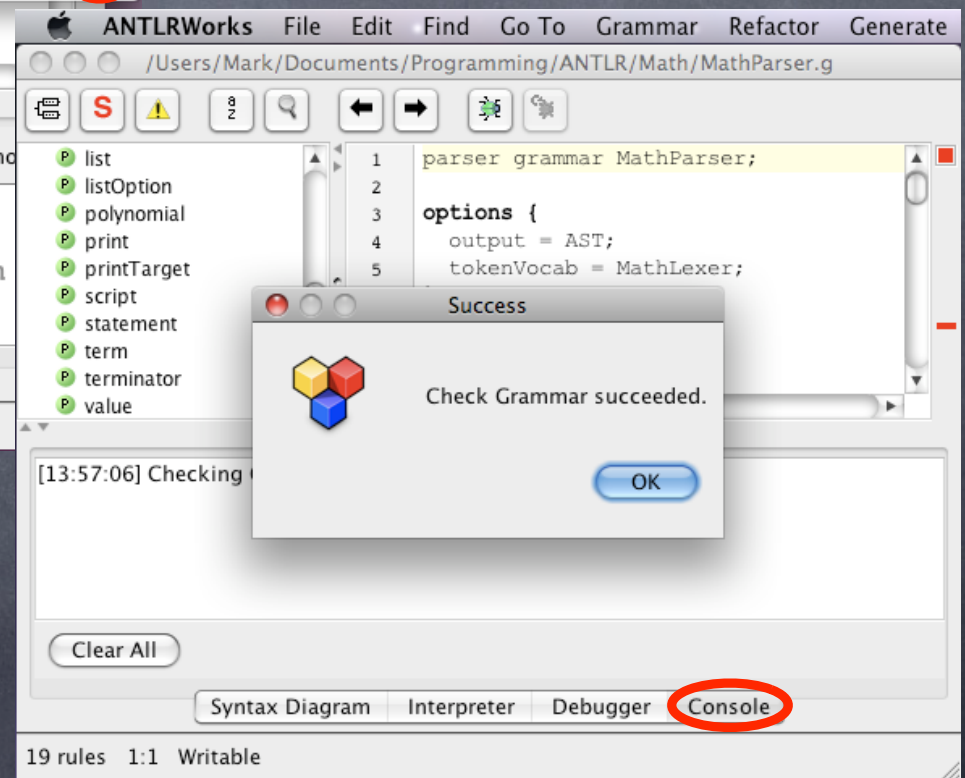
parser rule
syntax diagram



ANTLRWorks ...



grammar check
result

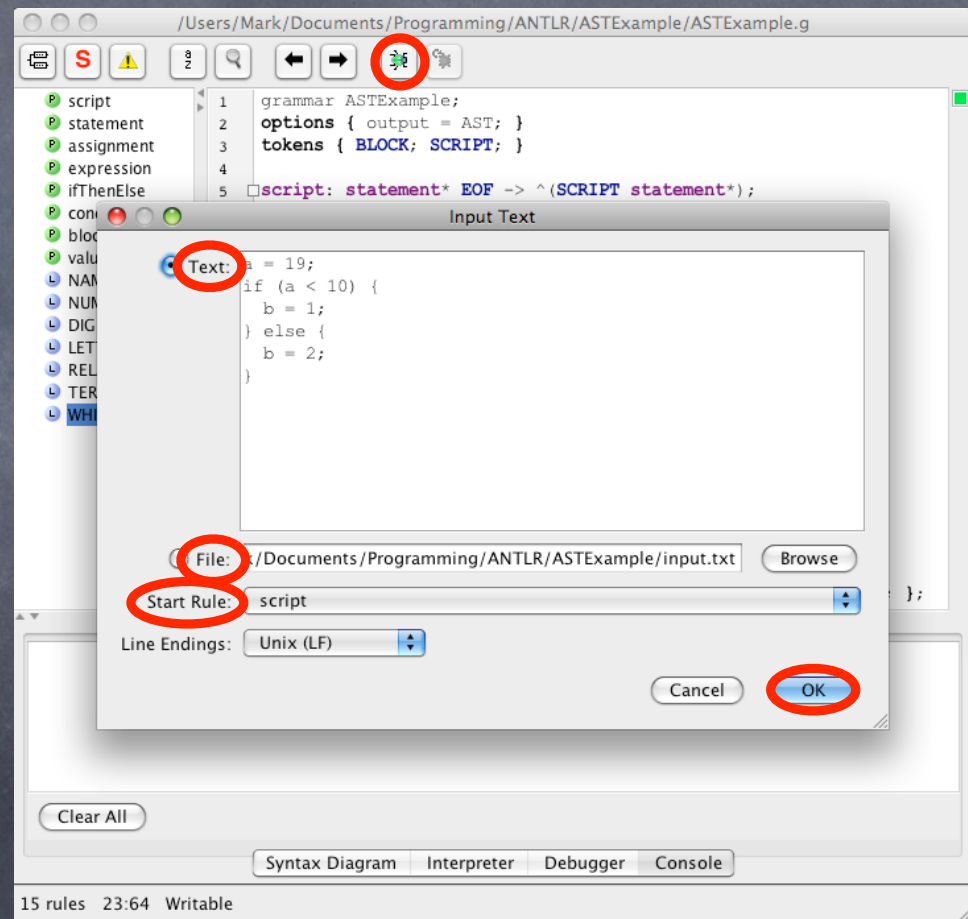


↑
requesting a
grammar check

ANTLRWorks Debugger

- Simple when lexer and parser rules are combined in a single grammar file

- press Debug toolbar button
- enter input text or select an input file
- select start rule
 - allows debugging a subset of grammar
- press OK button



ANTLRWorks Debugger ...

At the bottom of the ANTLRWorks window

The screenshot displays the ANTLRWorks Debugger interface. At the top, there is a toolbar with a stop button (circled in red) and a play button (circled in red). Below the toolbar, the 'Break on' section has checkboxes for 'All', 'Location', 'Consume' (checked), 'LT', and 'Exception'. The 'Break on Terminate' checkbox is also present.

The main area is divided into several panels:

- Input:** Contains the source code:

```
a = 19;  
if (a < 10) {  
    b = 1;  
} else {  
    b = 2;  
}
```
- AST:** Shows the Abstract Syntax Tree structure. The root is 'r0', which is a 'script' rule. It has two children: an '=' node and an 'if' node. The '=' node has children 'a' and '19'. The 'if' node has children '<', 'BLOCK', and 'BLOCK'. The first 'BLOCK' has children 'a' and '10'. The second 'BLOCK' has children '=' and '='. The first '=' has children 'b' and '1'. The second '=' has children 'b' and '2'.
- Events:** A list of events with their locations and descriptions:

#	Event
474	Consume [/<-1>,0:-1,@-1]
475	Location (6,25)
476	Nil node 41
477	Location (6,28)
478	Location (6,28)
479	Add child 7 to 41
480	Location (6,28)
481	Add child 40 to 41
482	LT -1 ()
483	Set token boundaries 41 (0,39)
484	Location (6,38)
485	Exit rule :script
486	Terminate

At the bottom, there is a toolbar with buttons for 'Input', 'Output', 'Parse Tree', 'AST', 'Stack', and 'Events'. The 'AST' button is circled in red. Below this toolbar, there are tabs for 'Syntax Diagram', 'Interpreter', 'Debugger' (circled in red), and 'Console'.

ASTExample
Start Rule: script

ANTLRWorks Debugger ...

- A bit more complicated when
lexer and parser rules are in separate files

See the ANTLR Wiki page
“When do I need to use remote debugging?” at
<http://www.antlr.org/wiki/pages/viewpage.action?pageId=5832732>

- We'll demonstrate this
after we see the Java code that
ties all the generated classes together
 - see slides 53-56

Using Rule Return Values

These are examples from our tree grammar.

The code in curly braces is a rule "action" written in the target language, in this case Java.

```
printTarget
: NUMBER { out($NUMBER); }
| STRING_LITERAL {
    String s = unescape($STRING_LITERAL.text);
    out(s.substring(1, s.length() - 1)); // remove quotes
}
| NAME { out(getVariable($NAME)); }
| ^(FUNCTION NAME) { out(getFunction(NAME)); }
| functionEval { out($functionEval.result); }
| derivative // handles own output
;
```

"unescape", "out",
"getFunction", "getVariable",
"evalFunction" and "toDouble"
are methods we wrote in
the tree grammar coming up.

```
functionEval returns [double result]
: ^(FUNCTION fn=NAME v=NUMBER) {
    $result = evalFunction($fn, toDouble($v));
}
| ^(FUNCTION fn=NAME v=NAME) {
    $result = evalFunction($fn, getVariable($v));
}
;
```

Rule Actions

- Add code before and/or after the generated code in the method generated for a rule
 - can be used for AOP-like wrapping of methods
- `@init { ... }`
 - inserts contained code before generated code
 - can be used to declare local variables used in actions of rule alternatives
 - used in our tree parser `polynomial` and `term` rules ahead
- `@after { ... }`
 - inserts contained code after generated code

Attribute Scopes

• Data is shared between rules in two ways

- passing parameters and/or returning values
- using attributes

same as options to share data between Java methods in the same class

• Attributes can be accessible to

- a single rule using `@init` to declare them
- a rule and all rules invoked by it - rule scope
- all rules that request the named global scope of the attributes

• Attribute scopes

- define collections of attributes that can be accessed by multiple rules
- two kinds, global and rule scopes

Attribute Scopes ...

Global scopes

- named scopes defined outside any rule

- define with

```
scope name {  
    type variable;  
    ...  
}
```

Use an @init rule action to initialize attributes.

- request access to the scope in a rule with

`scope name;`

To access multiple scopes, list them separated by spaces.

- rule actions access variables in the scope with
`$name::variable`

Rule scopes

- unnamed scopes defined inside a rule

- define with

```
scope {  
    type variable;  
    ...  
}
```

- rule actions in the defining rule and rules invoked by it access attributes in the scope with
`$rule-name::variable`

Our Tree Grammar

```
tree grammar MathTree;
```

```
options {  
    ASTLabelType = CommonTree;  
    tokenVocab = MathParser;  
}
```

We're going to process an AST whose nodes are of type CommonTree.

We're going to use the tokens defined in both our MathLexer and MathParser grammars. The MathParser grammar already includes the tokens defined in the MathLexer grammar.

```
@header {  
    package com.ociweb.math;  
  
    import java.util.Map;  
    import java.util.TreeMap;  
}
```

We want the generated parser class to be in this package.

```
@members {  
    private Map<String, Function> functionMap = new TreeMap<String, Function>();  
    private Map<String, Double> variableMap = new TreeMap<String, Double>();
```

We're using TreeMaps so the entries are sorted on their keys which is desired when listing them.

Our Tree Grammar ...

```
private void define(Function function) {  
    functionMap.put(function.getName(), function);  
}
```

This adds a Function
to our function Map.

```
private Function getFunction(CommonTree nameNode) {  
    String name = nameNode.getText();  
    Function function = functionMap.get(name);  
    if (function == null) {  
        String msg = "The function \"" + name + "\" is not defined.";  
        throw new RuntimeException(msg);  
    }  
    return function;  
}
```

This retrieves a Function
from our function Map
whose name matches the text
of a given AST tree node.

```
private double evalFunction(CommonTree nameNode, double value) {  
    return getFunction(nameNode).getValue(value);  
}
```

This evaluates a function
whose name matches the text
of a given AST tree node
for a given value.

Our Tree Grammar ...

```
private double getVariable(CommonTree nameNode) {  
    String name = nameNode.getText();  
    Double value = variableMap.get(name);  
    if (value == null) {  
        String msg = "The variable \"" + name + "\" is not set.";  
        throw new RuntimeException(msg);  
    }  
    return value;  
}
```

This retrieves the value of a variable from our variable Map whose name matches the text of a given AST tree node.

```
private static void out(Object obj) {  
    System.out.print(obj);  
}  
  
private static void outln(Object obj) {  
    System.out.println(obj);  
}
```

These just shorten the code for print and println calls.

Our Tree Grammar ...

```
private double toDouble(CommonTree node) {  
    double value = 0.0;  
    String text = node.getText();  
    try {  
        value = Double.parseDouble(text);  
    } catch (NumberFormatException e) {  
        throw new RuntimeException("Cannot convert \"" + text + "\" to a double.");  
    }  
    return value;  
}
```

This converts the text of a given AST node to a double.

```
private static String unescape(String text) {  
    return text.replaceAll("\\\\n", "\\n");  
}
```

```
} // @members
```

This replaces all escaped newline characters in a String with unescaped newline characters. It is used to allow newline characters to be placed in literal Strings that are passed to the print command.

Our Tree Grammar ...

```
script: statement*;
```

```
statement: assign | combine | define | interactiveStatement | print;
```

```
interactiveStatement: help | list;
```

```
assign: ^(ASSIGN NAME v=value) { variableMap.put($NAME.text, $v.result); };
```

This adds a variable to the variable map.

could also use \$value here

```
value returns [double result]
```

```
: NUMBER { $result = toDouble($NUMBER); }  
| NAME { $result = getVariable($NAME); }  
| functionEval { $result = $functionEval.result; }  
;
```

This returns a value as a double. The value can be a number, a variable name or a function evaluation.

```
functionEval returns [double result]
```

```
: ^(FUNCTION fn=NAME v=NUMBER) {  
    $result = evalFunction($fn, toDouble($v));  
}  
| ^(FUNCTION fn=NAME v=NAME) {  
    $result = evalFunction($fn, getVariable($v));  
}  
;
```

This returns the result of a function evaluation as a double.

Our Tree Grammar ...

define

```
: ^(DEFINE name=NAME variable=NAME polynomial) {  
    define(new Function($name.text, $variable.text, $polynomial.result));  
}  
;
```

This builds a Function object and adds it to the function map.

polynomial returns [Polynomial **result**]

scope { Polynomial current; } ←

This builds a Polynomial object and returns it.

```
@init { $polynomial::current = new Polynomial(); }  
: ^(POLYNOMIAL term[""] (s=SIGN term[$s.text]))* {  
    $result = $polynomial::current;  
}  
;
```

The "current" attribute in this rule scope is visible to rules invoked by this one, such as term.

There can be no sign in front of the first term, so "" is passed to the term rule.
The coefficient of the first term can be negative.
The sign between terms is passed to subsequent invocations of the term rule.

Our Tree Grammar ...

```
term[String sign]
@init { boolean negate = "-".equals(sign); }
: ^(TERM coefficient=NUMBER) {
    double c = toDouble($coefficient);
    if (negate) c = -c; // applies sign to coefficient
    $polynomial::current.addTerm(new Term(c));
}
| ^(TERM coefficient=NUMBER? variable=NAME exponent=NUMBER?) {
    double c = coefficient == null ? 1.0 : toDouble($coefficient);
    if (negate) c = -c; // applies sign to coefficient
    double exp = exponent == null ? 1.0 : toDouble($exponent);
    $polynomial::current.addTerm(new Term(c, $variable.text, exp));
}
;
```

This builds a Term object and adds it to the current Polynomial.

Our Tree Grammar ...

help

```
: HELP {  
    outln("In the help below");  
    outln("* fn stands for function name");  
    outln("* n stands for a number");  
    outln("* v stands for variable");  
    outln("");  
    outln("To define");  
    outln("* a variable: v = n");  
    outln("* a function from a polynomial: fn(v) = polynomial-terms");  
    outln("  (for example, f(x) = 3x^2 - 4x + 1)");  
    outln("* a function from adding or subtracting two others: " +  
        "fn3 = fn1 +|- fn2");  
    outln("  (for example, h = f + g)");  
    outln("");  
    outln("To print");  
    // some lines omitted for space  
    outln("To exit: exit or quit");  
}  
;
```

This outputs help
on our language
which is useful in
interactive mode.

Our Tree Grammar ...

`list`

```
: ^(LIST FUNCTIONS) {  
    println("# of functions defined: " + functionMap.size());  
    for (Function function : functionMap.values()) {  
        println(function);  
    }  
}  
  
| ^(LIST VARIABLES) {  
    println("# of variables defined: " + variableMap.size());  
    for (String name : variableMap.keySet()) {  
        double value = variableMap.get(name);  
        println(name + " = " + value);  
    }  
}  
;
```

This lists all the functions or variables that are currently defined.

Our Tree Grammar ...

`combine`

```
: ^(COMBINE fn1=NAME op=SIGN fn2=NAME fn3=NAME) {  
    Function f2 = getFunction(fn2);  
    Function f3 = getFunction(fn3);  
    if ("+" .equals($op.text)) {  
        define(f2.add($fn1.text, f3));  
    } else if ("-".equals($op.text)) {  
        define(f2.subtract($fn1.text, f3));  
    } else {  
        // This should never happen since SIGN is defined to be either "+" or "-".  
        throw new RuntimeException(  
            "The operator \"" + $op +  
            " cannot be used for combining functions.");  
    }  
}  
;
```

This adds or subtracts two functions to create a new one.

"\$fn1.text" is the name of the new function to create.

Our Tree Grammar ...

print

```
: ^(PRINT printTarget*)  
  { System.out.println(); };
```

This prints a list of printTargets then prints a newline.

printTarget

```
: NUMBER { out($NUMBER); }  
| STRING_LITERAL {  
    String s = unescape($STRING_LITERAL.text);  
    out(s.substring(1, s.length() - 1)); /* removes quotes */  
}  
| NAME { out(getVariable($NAME)); }  
| ^(FUNCTION NAME) { out(getFunction($NAME)); }  
| functionEval { out($functionEval.result); }  
| derivative  
;
```

This prints a single printTarget without a newline.

on slide 46

derivative

```
: ^(DERIVATIVE NAME) {  
    out(getFunction($NAME).getDerivative());  
}  
;
```

This prints the derivative of a function.
This also could have been done
in place in the printTarget rule.

Using Generated Classes

• Our manually written Processor class

• uses the generated classes

- MathLexer extends Lexer
- MathParser extends Parser
- MathTree extends TreeParser

Lexer, Parser and TreeParser
extend BaseRecognizer

• uses other manually written classes

- Function
- Polynomial
- Term

• supports two modes

- batch – see processFile method
- interactive – see processInteractive method

Processor.java

```
package com.ociweb.math;

import java.io.*;
import java.util.Scanner;
import org.antlr.runtime.*;
import org.antlr.runtime.tree.*;

public class Processor {

    public static void main(String[] args) throws IOException, RecognitionException {
        if (args.length == 0) {
            new Processor().processInteractive();
        } else if (args.length == 1) { // name of file to process was passed in
            new Processor().processFile(args[0]);
        } else { // more than one command-line argument
            System.err.println("usage: java com.ociweb.math.Processor [file-name]");
        }
    }
}
```

Processor.java ...

```
private void processFile(String filePath) throws IOException, RecognitionException {
    CommonTree ast = getAST(new FileReader(filePath));
    //System.out.println(ast.toStringTree()); // for debugging
    processAST(ast);
}

private CommonTree getAST(Reader reader) throws IOException, RecognitionException {
    MathParser tokenParser = new MathParser(getTokenStream(reader));
    MathParser.script_return parserResult = tokenParser.script(); // start rule method
    reader.close();
    return (CommonTree) parserResult.getTree();
}

private CommonTokenStream getTokenStream(Reader reader) throws IOException {
    MathLexer lexer = new MathLexer(new ANTLRReaderStream(reader));
    return new CommonTokenStream(lexer);
}

private void processAST(CommonTree ast) throws RecognitionException {
    MathTree treeParser = new MathTree(new CommonTreeNodeStream(ast));
    treeParser.script(); // start rule method
}
```

Processor.java ...

```
private void processInteractive() throws IOException, RecognitionException {  
    MathTree treeParser = new MathTree(null); // a TreeNodeStream will be assigned later  
    Scanner scanner = new Scanner(System.in);  
  
    while (true) {  
        System.out.print("math> ");  
        String line = scanner.nextLine().trim();  
        if ("quit".equals(line) || "exit".equals(line)) break;  
        processLine(treeParser, line);  
    }  
}
```

Processor.java ...

```
private void processLine(MathTree treeParser, String line) throws RecognitionException {  
    // Run the lexer and token parser on the line.  
    MathLexer lexer = new MathLexer(new ANTLRStringStream(line));  
    MathParser tokenParser = new MathParser(new CommonTokenStream(lexer));  
    MathParser.statement_return parserResult = tokenParser.statement(); // start rule method  
  
    // Use the token parser to retrieve the AST.  
    CommonTree ast = (CommonTree) parserResult.getTree();  
    if (ast == null) return; // line is empty  
  
    // Use the tree parser to process the AST.  
    treeParser.setTreeNodeStream(new CommonTreeNodeStream(ast));  
    treeParser.statement(); // start rule method  
}  
  
} // end of Processor class
```

We can't create a new instance of MathTree for each line processed because it maintains the variable and function Maps.

ANTLRWorks Debugger

- Let's demonstrate using remote debugging which is necessary when lexer and parser rules are in separate grammar files
 - edit `build.properties` to include `-debug in tool.options`
 - `ant clean run`
 - the run target in `build.xml` tells it to parse the file "simple.math"
 - start ANTLRWorks
 - open the parser grammar file
 - select Debugger ... Debug Remote...
 - press "Connect" button
 - debug as usual

ANTLRWorks Debugger ...

The screenshot displays the ANTLRWorks Debugger interface. The top panel shows the `MathParser.g` grammar file with the following content:

```
57 print
58 : PRINT printTarget* terminator -> ^(PRINT printTarget*);
59
60 printTarget
61 : NUMBER -> NUMBER
62 | s1=STRING_LITERAL -> $s1
63 | NAME -> NAME
64 // This is a function reference to print a string representation.
65 | NAME LEFT_PAREN RIGHT_PAREN -> ^(FUNCTION NAME)
66 | functionEval
67 | derivative
68 ;
69
70 // This is the "start rule".
71 script: statement* EOF -> ^(SCRIPT statement*)
```

The bottom panel shows the debugger's state. The **Input** tab is selected, displaying the following text:

```
a = 3.14
f(x) = 3x^2 - 4x + 2
print "The value of f for " a " is " f(a)
print "The derivative of " f() " is " f'()
list variables
list functions
g(y) = 2y^3 + 6y - 5
h = f + g
print h()
```

The **AST** tab is also selected, showing the Abstract Syntax Tree for the input. The root node is `r0`, which is a `MathParser.g` rule. The tree structure is as follows:

- `r0` (Root)
 - `=` (Assignment)
 - `a` (Variable)
 - `3.14` (Number)
 - `DEFINE` (Function Definition)
 - `f` (Function Name)
 - `x` (Parameter)
 - `POLYNOMIAL` (Polynomial)
 - `TERM` (Term)
 - `3` (Coefficient)
 - `x` (Variable)
 - `2` (Exponent)
 - `-` (Operator)
 - `TERM` (Term)
 - `4` (Coefficient)
 - `x` (Variable)
 - `+` (Operator)
 - `TERM` (Term)
 - `2` (Constant)

The **Input** and **AST** tabs are circled in red. The **Debugger** tab is also visible at the bottom.

References

- ANTLR

- <http://www.antlr.org>

- ANTLRWorks

- <http://www.antlr.org/works>

- StringTemplate

- <http://www.stringtemplate.org>

- http://www.codegeneration.net/tiki-read_article.php?articleId=65 and 77

- My slides and code examples

- <http://www.ociweb.com/mark> - look for "ANTLR 3"